## Optimal real-time transmission of Markov sources under constraints on the number of transmissions

#### Aditya Mahajan Joint work with Jhelum Chakravorty

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Information Theory Seminar, University of Toronto 14 Nov, 2014

- Sequential transmission of data
- > Zero- (or finite-) delay reconstruction



#### Sensor Networks

Sequential transmission of data Zero- (or finite-) delay reconstruction







#### Smart Grids

Sequential transmission of data Zero- (or finite-) delay reconstruction







#### Internet of Things

Real-time transmission of Markov sources- (Aditya Mahajan)

Sequential transmission of data Zero- (or finite-) delay reconstruction





#### Internet of Things







Real-time transmission of Markov sources- (Aditya Mahajan)

## The communication system



Source  $\blacktriangleright X_t \in \mathbb{Z}$ 

> First-order time-homogeneous symmetric Markov source.

$$\begin{array}{ll} \mbox{Transmitter} & U_t = f_t(X_{1:t}, U_{1:t-1}) \mbox{ and } Y_t = \begin{cases} X_t & \mbox{if } U_t = 1 \\ \epsilon & \mbox{if } U_t = 0 \end{cases} \end{array}$$

$$\begin{array}{l} \text{Receiver} & \blacktriangleright \ \hat{X}_t = g_t(Y_{1:t}) \\ & \blacktriangleright \ \text{Distortion:} \ d(X_t - \hat{X}_t) \ \text{where} \ d(e) = d(-e) \leqslant d(e+1) \end{array}$$

 $\label{eq:communication} \begin{array}{l} \bullet \mbox{ Transmission strategy } f = \{f_t\}_{t=0}^{\infty}. \\ \hline \mbox{ Strategies } \bullet \mbox{ Estimation strategy } g = \{g_t\}_{t=0}^{\infty}. \end{array}$ 



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#### The constrained optimization problem

$$\min_{(f,g)} D_{\beta}(f,g) \quad \text{ such that } N_{\beta}(f,g) \leqslant \alpha$$

Minimize expected distortion such that expected # of transmissions is less than  $\alpha$ 

iscounted  
setup  
$$D_{\beta}(f,g) = (1-\beta) \mathbb{E}^{(f,g)} \left[ \sum_{t=0}^{\infty} \beta^{t} d(X_{t} - \hat{X}_{t}) \middle| X_{0} = 0 \right]$$
$$N_{\beta}(f,g) = (1-\beta) \mathbb{E}^{(f,g)} \left[ \sum_{t=0}^{\infty} \beta^{t} U_{t} \middle| X_{0} = 0 \right]$$

Average cost setup

D

$$D_{1}(f,g) = \limsup_{T \to \infty} \frac{1}{T} \Big[ \sum_{t=0}^{T-1} d(X_{t} - \hat{X}_{t}) \ \Big| \ X_{0} = 0$$
$$N_{1}(f,g) = \limsup_{T \to \infty} \frac{1}{T} \Big[ \sum_{t=0}^{T-1} U_{t} \ \Big| \ X_{0} = 0 \Big]$$





#### Assumptions on the model

(Ao)  $X_t \in \mathbb{Z}$ , and  $X_0 = 0$ .

(A1) The transition matrix is Toeplitz with decaying off-diagonal terms.

$$P = \begin{bmatrix} \ddots & p_0 & \ddots & & \\ \cdots & p_1 & p_0 & p_1 & \cdots & \\ & \ddots & p_1 & p_0 & p_1 & \cdots \\ & & \ddots & \ddots & p_0 & \ddots \end{bmatrix} \text{ and } \begin{array}{c} p_0 \geqslant p_1 \geqslant p_2 \geqslant \cdots \\ p_0 > 0 \end{array}$$

▶ Nayyar et al, assumed that the transistion matrix was banded, that is,  $\exists b$  such that  $p_k = 0$ , for all  $k \ge b$ .

(A2) The distortion function is even and increasing on  $\mathbb{Z}_{\geq 0}$ .

 $\forall e \in \mathbb{Z}_{\geqslant 0}: \quad d(e) = d(-e) \quad \text{and} \quad d(e) \leqslant d(e+1).$ 

Furthermore,

 $\mathbf{d}(\mathbf{0}) = \mathbf{0} \quad \text{and} \quad \mathbf{d}(e) \neq \mathbf{0}, \quad \forall e \neq \mathbf{0}.$ 

#### An example: Symmetric birth-death Markov Chain

$$P_{ij} = \begin{cases} p, & \text{if } |i-j| = 1; \\ 1-2p, & \text{if } i = j; \\ 0, & \text{otherwise,} \end{cases} \quad \text{where } p \in (0, \frac{1}{2}), \quad d(e) = |e|$$





#### Main results: Distortion-transmission function

Distortiontransmission function

$$D^*_{\beta}(\alpha) = \min\{D_{\beta}(f,g) \text{ such that } N_{\beta}(f,g) < \alpha\}$$

Properties:

- $D^*_{\beta}(\alpha)$  is convex and decreasing.
- $\blacktriangleright \lim_{\alpha \to 0} D^*_\beta(\alpha) = \infty \text{ and } \lim_{\alpha \to 1} D^*_\beta(\alpha) = 0$



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#### Main results: Optimal communication strategies

EstimationLet Zt be the most recently transmitted symbol up to time t. Then, thestrategyoptimal transmission strategy is

 $g^*(Y_{1:t}) = Z_t.$ 

 $\begin{array}{ll} \mbox{Transmission} & \mbox{Let } E_t = X_t - Z_t \mbox{ and } f^{(k)} \mbox{ be a threshold-based strategy given by} \\ & \mbox{strategy} \\ & f^{(k)}(X_t,Y_{1:t-1}) = \begin{cases} 0, & \mbox{if } |E_t| \leqslant k, \\ 1, & \mbox{if } |E_t| > k. \end{cases} \end{array}$ 

The optimal transmission strategy is a possibly randomized strategy that, at each stage picks

- $f^{(k^*)}$  with probability probability  $\theta^*$
- $f^{(k^*+1)}$  with probability probability  $1 \theta^*$
- Let  $N_{\beta}^{(k)} = N_{\beta}(f^{(k)}, g^*)$  and  $D_{\beta}^{(k)} = D_{\beta}(f^{(k)}, g^*)$ . Then:
- $\blacktriangleright k^*$  is the largest k such that  $N_\beta^{(k)} \geqslant \alpha)$
- $\theta^*$  is such that

$$\theta^* N_\beta^{(\mathbf{k}^*)} + (1-\theta^*) N_\beta^{(\mathbf{k}^*+1)} = \alpha$$



### Main results: Optimal communication strategies





 $\theta^* N_{\beta}^{(\mathbf{k}^*)} + (1 - \theta^*) N_{\beta}^{(\mathbf{k}^* + 1)} = \alpha$ 

Real-time transmission of Markov sources- (Aditya Mahajan)

Standard > Achievability: Identify a good strategy and evaluate its performance.
 technique > Converse: Determine a lower bound on distortion.



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Our approach Model the optimization problem as a decentralized stochastic control problem. [Witsenhausen 1979, Walrand-Varaiya 1983, Teneketzis 2006, Mahajan-Teneketzis 2009, ...]

▶ The system has two decision makers: the transmitter and the estimator, that have access to different information.



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- ▶ The system has two decision makers: the transmitter and the estimator, that have access to different information.
- Identify qualitative properties of optimal strategies
- Identify a dynamic programming decomposition
- Determine optimal strategies based on the dynamic program.



## Outline of the proof: Setting up a countable state DP

Step 1: Identify an information state and dynamic program for Lagrange relaxation Use the common-information approach of Nayyar-Mahajan-Teneketzis 2013 to transform the decentralized control problem to a centralized coordination problem

Step 2: Determine qualitative properties of optimal strategies from the DP Use majorization theory to show that, under an optimal policy, the reachable set of the information state is an almost symmetric and unimodal distribution.

The optimal transmission strategy is of a threshold-type and the optimal estimation strategy does not depend on the value of the threshold.

Previously proved by [Lipsa-Marins 2011, Nayyar-Başar-Teneketzis-Veeravalli 2013]

#### Step 3: Fix the estimator. Investigate the best-response transmitter

Use standard results from DP to identify sufficient conditions under which optimal strategy is time-homogeneous and given by the unique fixed-point of a DP.



## Outline of the proof: Identifying a solution to the DP

Step 4: Evaluate cost of Lagrange relaxation for a particular transmission strategy Both estimation and transmission strategies are fixed. Solve the DP to obtain renewaltheory-like relationships.

Step 5: Identify Lagrange multipliers for which a particular strategy is optimal Similar to the idea of calibration in multi-armed bandits.

# Step 6: Evaluate optimal Lagrange performance. Infer the optimal strategy for the constrained setup

The optimal Lagrange performance is continuous, piecewise linear, concave, and increasing in the Lagrange multiplier.

Show that a Bernoulli randomized simple transmission strategy is optimal.

The performance of the optimal strategy gives the distortion-transmission function.

#### Lagrange Relaxation



Minimize expected distortion such that expected # of transmissions is less than  $\boldsymbol{\alpha}$ 



#### Lagrange Relaxation

$$\min_{(\mathfrak{f},g)} D_{\beta}(\mathfrak{f},g) \quad \text{ such that } N_{\beta}(\mathfrak{f},g) \leqslant \alpha$$

Minimize expected distortion such that expected # of transmissions is less than  $\boldsymbol{\alpha}$ 

$$\begin{array}{ll} \mbox{Lagrange} & C^*_\beta(\lambda)\coloneqq \inf_{(f,g)} C_\beta(f,g;\lambda) & \mbox{where } C_\beta(f,g;\lambda) = D_\beta(f,g) + \lambda N_\beta(f,g) \\ \mbox{Relaxation} & \end{array}$$



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Search space of strategies (f, g)

- Restrict the search space of strategies (f, g) by identifying structure of optimal tranmission and estimation strategies.
- Difficulty: Non-classical information structure



#### Step 1a: Removing irrelevant information

 $\label{eq:Information} \begin{array}{ll} I_t = \{X_{1:t}, U_{1:t-1}, Y_{1:t-1}\} \text{ and } J_t = \{Y_{1:t}\}. \\ structure \end{array}$ 



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**Remove** > Aritarily fix estimation strategy.

irrelevant data 
Finding the best-response transmitter is a centralized stochastic control problem.

►  $\tilde{I}_t = \{X_t, Y_{1:t-1}\}$  is a controlled Markov process.  $\mathbb{P}(\tilde{I}_{t+1} | I_t, U_t, Y_t) = \mathbb{P}(\tilde{I}_{t+1} | \tilde{I}_t, Y_t);$   $\mathbb{E}[d(X_t - \hat{X}_t) + \lambda U_t | I_t, U_t, Y_t] = \mathbb{E}[d(X_t - \hat{X}_t) + \lambda U_t | \tilde{I}_t, Y_t]$ 

 $\blacktriangleright$  Therefore, there is no loss of optimality in using control  $U_t = \tilde{f}_t(\tilde{I}_t).$ 

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 $\blacktriangleright$  Therefore, there is no loss of optimality in using control  $U_t = \tilde{f}_t(\tilde{I}_t).$ 

 $\label{eq:simplified} \tilde{I}_t = \{X_t, Y_{1:t-1}\} \text{ and } J_t = \{Y_{1:t}\}.$  Info Struct



#### Step 1b: Equivalent centralized problem

Info Struct  $\tilde{I}_t = \{X_t, Y_{1:t-1}\}$  and  $J_t = \{Y_{1:t}\}$ .

$$X_{t-1}, Y_{1:t-2}$$
 $Y_{1:t-1}$ 
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<sup>▶</sup> Nayyar, Mahajan and Teneketzis, "Decentralized stochastic control with partial history sharing: A common information approach," IEEE TAC 2013.

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$$X_{t-1}, Y_{1:t-2}$$
  $Y_{1:t-1}$   $X_t, Y_{1:t-1}$   $Y_{1:t}$ 



Coordinated system is equivalent to original system

Nayyar, Mahajan and Teneketzis, "Decentralized stochastic control with partial history sharing: A common information approach," IEEE TAC 2013.



#### Step 1c: Structural results and dynamic program

 $\begin{array}{ll} \mbox{Information} & \blacktriangleright \mbox{Pre-transmission belief: } \Pi_t(x) = \mathbb{P}(X_t = x \mid Y_{1:t-1}) \\ & \mbox{states} & \blacktriangleright \mbox{Post-transmission belief: } \Phi_t(x) = \mathbb{P}(X_t = x \mid Y_{1:t}). \end{array}$ 



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 $\label{eq:construct} \begin{array}{ll} \mbox{Structural} & \mbox{There is no loss of optimality in using} \\ \mbox{results} & \mbox{$U_t=f_t(\Pi_t)$, and $\hat{X}_t=g_t(\Phi_t)$.} \end{array}$ 

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Dynamic program

 $W_{\mathsf{T}+1}(\pi) = 0$ 

and for  $t = T, \ldots, 1$ 

$$\begin{split} V_t(\phi) &= \min_{\hat{x} \in \mathcal{X}} \mathbb{E}[d(X_t, \hat{x}) + W_{t+1}(\Pi_{t+1}) \mid \Phi_t = \phi] \\ W_t(\pi) &= \min_{\mathbf{T}} \mathbb{E}[\lambda U_t + V_t(\Phi_t) \mid \Pi_t = \pi] \end{split}$$



Can we use the DP to say something more about the optimal strategy?

### Qualitative properties of optimal stategies

#### [Imer-Başar 2005 & 2010]

Fixed number of transmissions for finite horizon LQG setup.

#### [Lipsa-Martins 2009 & 2011, Molin-Hirche 2009]

Remote estimation with communication cost for finite horizon LQG setup.

#### [Nayyar-Başar-Teneketzis-Veeravalli 2013]

Remote estimation with communication cost for finite horizon Markov chain setup. Also considered energy harvesting at the transmitter.


#### Step 2a: a.s.u. distributions and majorization

a.s.u. distribution A probability distribution  $\mu$  over  $\mathbb{Z}$  is said to be almost summetric and unimodal about a pont  $\alpha$  if

 $\mu_{a+n} \geqslant \mu_{a-n} \geqslant \mu_{a+k+1}.$ 

a.s.u. The a.s.u. rearrangement of a probability distribution  $\mu$ , denoted by  $\mu^+$  rearrangement is a permutation of  $\mu$  such that for every n

 $\mu_n^+ \geqslant \mu_n^- \geqslant \mu_{n+1}^+$ 

 $\begin{array}{ll} \mbox{Majorization} & \mbox{A probability distribution } \mu \mbox{ majorizes a distribution } \nu, \mbox{ denoted by } \mu \succeq_m \\ \nu \mbox{ if for all } n \end{array}$ 

$$\sum_{i=-k}^k \mu_i^+ \geqslant \sum_{i=-k}^k \nu_i^+$$



# Step 2b: Qualitative properties of optimal strategies

#### Monotonicity of value functions

If  $\tilde{\phi}$  is an a.s.u. distrution such that  $\tilde{\phi} \succeq_m \phi$ , then  $V_t(\phi) \ge V_t(\tilde{\phi})$ .

#### Structure of optimal estimator

If  $\phi_t$  is a.s.u. about a, then the optimal estimate is a.

#### Structure of optimal transmitter

If  $\pi_t$  is a.s.u. about a, then the optimal prescription  $\gamma_t$  is of the form

$$\gamma_t(x) = \begin{cases} 1, & |x-a| \geqslant k(\pi_t) \\ 0, & |x-a| < k(\pi_t) \end{cases}$$

Real-time transmission of Markov sources- (Aditya Mahajan)



Nayyar, Başar, Teneketzis, Veeravalli, TAC 2013

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Using these properties, one can show that under an optimal strategy  $\pi_t$  and  $\phi_t$  are a.s.u.

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Nayyar, Başar, Teneketzis, Veeravalli, TAC 2013

# Step 2c: Structure of optimal estimator (Nayyar et al, 2013)

$$Z_0 = 0 \quad \text{and} \quad Z_t = \begin{cases} X_t & \text{if } U_t = 1; \\ Z_{t-1} & \text{if } U_t = 0. \end{cases}$$

The estimator can keep track of  $Z_t$  as follows:

$$Z_0 = 0 \quad \text{and} \quad Z_t = \begin{cases} Y_t & \text{if } Y_t \neq \epsilon; \\ Z_{t-1} & \text{if } Y_t = \epsilon. \end{cases}$$

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Theorem 1 The process  $\{Z_t\}_{t=0}^{\infty}$  is a sufficient statistic at the estimator and an optimal estimation strategy is given by  $\hat{X}_t = g_t^*(Z_t) = Z_t$  (\*)

**Remark** > The optimal estimation strategy is time-homogeneous and can be specified in closed form.

# **Step 2d:** Structure of optimal transmitter (Nayyar et al)

Error process Let  $E_t=X_t-Z_{t-1}$  denote the error process.  $\{E_t\}_{t=0}^\infty$  is a controlled Markov process where

$$E_0 = 0 \quad \text{and} \quad \mathbb{P}(E_{t+1} = n \mid E_t = e, U_t = u) = \begin{cases} P_{0n}, & \text{if } u = 1; \\ P_{en}, & \text{if } u = 0. \end{cases}$$



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**Theorem 2** When the estimation strategy is of the form (\*), then  $\{E_t\}_{t=0}^{\infty}$  is a sufficient statistic at the transmitter.

Furthermore, an optimal transmission strategy is characterized by a time-varying threshold  $\{k_t\}_{t=0}^\infty,$  i.e.,

$$U_t = f_t(E_t) = \begin{cases} 1 & \text{if } |E_t| \ge k_t; \\ 0 & \text{if } |E_t| < k_t. \end{cases}$$



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- Proof idea ► The proof of [Nayyar et al, 2013] was based on some majorization inequalities of [Hajek et al, 2009] for distributions with finite support.
  - ► We extend these inequalities to distributions over integers using results of [Wang-Woo-Madiman, 2014].





We have identified the structure of optimal transmission and estimation strategies for the finite-horizon Lagrange relaxation of the original problem.

How do these results extend to infinite horizon setup?

# **Step 3:** Infinite horizon setup (for Lagrange relaxation)

Main idea • Based on Thm 1, restrict attention to time-homogeneous estimation strategy

 $\widehat{X}_t = \mathbf{g}_t^*(Z_t) = Z_t$ 

 Consider the problem of finding the "best response" estimation strategy.



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- Standard MDP results apply under mild technical assumptions.

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- Standard MDP results apply under mild technical assumptions.

Assum (A3) For every  $\lambda \ge 0$ , there exists a function  $w : \mathbb{Z} \to \mathbb{R}$  and postive and finite constants  $\mu_1$  and  $\mu_2$  such that for all  $e \in \mathbb{Z}$ , we have that

 $\text{max}\{\lambda,d(e)\}\leqslant \mu_1w(e)$ 

$$\max\Big\{\sum_{n=-\infty}^{\infty}\mathsf{P}_{en}w(n),\sum_{n=-\infty}^{\infty}\mathsf{P}_{0n}w(n)\Big\}\leqslant \mu_2w(e).$$



# Step 3: Structure of optimal transmitter for infinite horizon

Structure Under assumption (A3), optimal transmission strategy is characterized by time-homogeneous threshold k, i.e.,

$$U_t = f(E_t) = \begin{cases} 1 & \text{if } |E_t| \ge k; \\ 0 & \text{if } |E_t| < k. \end{cases}$$



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Dynamic For  $\beta \in (0, 1)$ , the optimal strategy is determined by the unique fixed program point of the following DP:

$$\begin{split} V_{\beta}(e;\lambda) &= \min \left\{ (1-\beta)\lambda + \beta \sum_{n=-\infty}^{\infty} \mathsf{P}_{0n} V_{\beta}(n;\lambda), \quad \begin{array}{l} \text{Transmit} \\ (1-\beta)d(e) + \beta \sum_{n=-\infty}^{\infty} \mathsf{P}_{en} V_{\beta}(n;\lambda) \right\} \quad \begin{array}{l} \text{Don't} \\ \text{Transmit} \\ \end{array} \end{split}$$



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$$U_t = f(E_t) = \begin{cases} 1 & \text{if } |E_t| \ge k; \\ 0 & \text{if } |E_t| < k. \end{cases}$$

**Dynamic** For  $\beta \in (0, 1)$ , the optimal strategy is determined by the unique fixed program point of the following DP:

$$\begin{split} V_{\beta}(e;\lambda) &= \min \left\{ (1-\beta)\lambda + \beta \sum_{n=-\infty}^{\infty} \mathsf{P}_{0n} V_{\beta}(n;\lambda), \quad \begin{array}{l} \text{Transmit} \\ (1-\beta)d(e) + \beta \sum_{n=-\infty}^{\infty} \mathsf{P}_{en} V_{\beta}(n;\lambda) \right\} \quad \begin{array}{l} \text{Don't} \\ \text{Transmit} \\ \end{array} \end{split}$$

 $\begin{array}{ll} \mbox{Lagrange} & \mbox{Let } f^*_\beta(\cdot;\lambda) \mbox{ be the time-homogeneous optimal transmission strategy.} \\ \mbox{relaxation} & C^*_\beta(\lambda) \coloneqq \inf_{(f,q)} C_\beta(f,g;\lambda) = C_\beta(f^*_\beta,g^*;\lambda) = V_\beta(0;\lambda) \end{array}$ 



### Step 3: The SEN Cond. and the long-term average setup

**SEN Conditions** For any  $\lambda \ge 0$ , the value function  $V_{\beta}(\cdot; \lambda)$  satisfy the SEN condition:

- $\label{eq:sigma} \mbox{(S1)} \quad \mbox{There exists a reference state $e_0 \in \mathbb{Z}$ such that $V_\beta(e_0;\lambda) < \infty$ for all $\beta \in (0,1)$.}$
- (S2) Define  $h_{\beta}(e;\lambda) = (1-\beta)^{-1}[V_{\beta}(e;\lambda) V_{\beta}(e_0;\lambda)]$ . There exists a function  $K_{\lambda} : \mathbb{Z} \to \mathbb{R}$  such that  $h_{\beta}(e;\lambda) \leq K_{\lambda}(e)$  for all  $e \in \mathbb{Z}$  and  $\beta \in (0,1)$ .
- (S3) There exists a non-negative (finite) constant  $L_{\lambda}$  such that  $-L_{\lambda} \leq h_{\beta}(e;\lambda)$  for all  $e \in \mathbb{Z}$  and  $\beta \in (0,1)$ .

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- (S3) There exists a non-negative (finite) constant  $L_{\lambda}$  such that  $-L_{\lambda} \leq h_{\beta}(e;\lambda)$  for all  $e \in \mathbb{Z}$  and  $\beta \in (0,1)$ .
- Proof ideas ► The Markov chain induced by f<sup>(0)</sup> is 0-standard, i.e., for every state e, the expected time and expected cost for first passage to 0 is finite.
  - Hence, (S1) and (S2) hold.
  - ▶ For any  $e \in \mathbb{Z}_{\geq 0}$ ,  $[P]_{e+1} >_r [P]_e$ , where  $>_r$  denotes reflected stochastic dominance.
  - Using induction show that for every  $\lambda$ ,  $V_{\beta}(e; \lambda)$  is even and increasing in *e*.
  - Hence (S3) holds.



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Vanishing Let  $f_1^*(\cdot; \lambda)$  be any limit point of  $f_{\beta}^*(\cdot; \lambda)$  as  $\beta \uparrow 1$ . discount Then the time-homogeneous transmission strategy  $f_1^*(\cdot; \lambda)$  is optimal approach for  $\beta = 1$  (the long-term average setup).

Furthermore, the performance of this optimal strategy is

$$C_1^*(\lambda) \coloneqq \inf_{(f,g)} C_1(f,g;\lambda) = C_1(f_1^*,g^*;\lambda) = \lim_{\beta \uparrow 1} V_\beta(0;\lambda) = \lim_{\beta \uparrow 1} C_\beta^*(\lambda).$$



Time-homogeneous threshold-based transmission strategies are optimal.

The optimal threshold can be determined by solving a DP.

The DP is well-behaved. The long-term average setup may be analyzed using the vanishing discount approach.

So what? Does the DP give any insights beyond numerical computations?

Threshold- We analyze the performace of  $(f^{(k)}, g^*)$ , where based strategy  $f^{(k)}(e) \coloneqq \begin{cases} 1, & \text{if } |e| \ge k; \\ 0, & \text{if } |e| < k. \end{cases}$ 

Threshold- We analyze the performace of  $(f^{(k)}, g^*)$ , where based strategy  $f^{(k)}(e) \coloneqq \begin{cases} 1, & \text{if } |e| \ge k; \\ 0, & \text{if } |e| < k. \end{cases}$ 

Performance of  $\blacktriangleright$  Distortion  $D_{\beta}^{(k)}(e)$  under strategy  $(f^{(k)}, g^*)$ : a given strategy  $\begin{pmatrix} \beta & \sum \\ -\beta & -p_{0n} D_{2n}^{(k)}(n) \end{pmatrix}$ .

$$D_{\beta}^{(k)}(e) = \begin{cases} \beta \sum_{n=-\infty}^{\infty} P_{0n} D_{\beta}^{(k)}(n), & |e| \ge k\\ (1-\beta)d(e) + \beta \sum_{n=-\infty}^{\infty} P_{en} D_{\beta}^{(k)}(n), & |e| < k \end{cases}$$

• Transmissions  $N_{\beta}^{(k)}(e)$  under strategy  $(f^{(k)}, g^*)$ :

$$N_{\beta}^{(k)}(e) = \begin{cases} (1-\beta) + \beta \sum_{n=-\infty}^{\infty} P_{0n} N_{\beta}^{(k)}(n), & |e| \ge k \\ \beta \sum_{n=-\infty}^{\infty} P_{en} N_{\beta}^{(k)}(n), & |e| < k \end{cases}$$



 $\begin{array}{ll} \mbox{Cost until first} & \mbox{Let } S^{(k)}=\{e\in\mathbb{Z}:|e|\leqslant k-1\}\mbox{ and }\tau^{(k)}\mbox{ be escape time of set }S^{(k)}.\\ & \mbox{transmission} \end{array}$ 



Cost until first Let  $S^{(k)} = \{e \in \mathbb{Z} : |e| \leq k-1\}$  and  $\tau^{(k)}$  be escape time of set  $S^{(k)}$ . transmission

$$\begin{array}{ll} \text{Define} \quad L_{\beta}^{(k)} \coloneqq \mathbb{E}\left[\left.\sum_{t=0}^{\tau^{(k)}-1}\beta^{t}d(E_{t})\middle|E_{0}=0\right]\right.\\ \\ \mathcal{M}_{\beta}^{(k)} \coloneqq \frac{1-\mathbb{E}[\beta^{\tau^{(k)}}\mid E_{0}=0]}{1-\beta} \end{array} \right.$$

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Renewal relationships

$$D_{\beta}^{(k)} = rac{L_{\beta}^{(k)}}{M_{\beta}^{(k)}}$$
 and  $N_{\beta}^{(k)} = rac{1}{M_{\beta}^{(k)}} - (1 - \beta)$ 

We show that these expressions satisfy the recursive relationships shown on the previous slide.

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Renewal relationships

$$\begin{split} D_{\beta}^{(k)} &= \frac{L_{\beta}^{(k)}}{M_{\beta}^{(k)}} \quad \text{and} \quad N_{\beta}^{(k)} &= \frac{1}{M_{\beta}^{(k)}} - (1 - \beta) \\ L_{1}^{(k)} &= \lim_{\beta \uparrow 1} L_{\beta}^{(k)}, \quad M_{1}^{(k)} &= \lim_{\beta \uparrow 1} M_{\beta}^{(k)}. \end{split}$$

Vanishing discount relationships

$$\begin{split} D_1^{(k)} &= \lim_{\beta \uparrow 1} D_\beta^{(k)} = \frac{L_1^{(k)}}{M_1^{(k)}} \\ N_1^{(k)} &= \lim_{\beta \uparrow 1} N_\beta^{(k)} = \frac{1}{M_1^{(k)}} \end{split}$$

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and



### Step 4: Computing performance

 $\begin{array}{ll} \mbox{Analytic} & \mbox{Let } P^{(k)} \mbox{ and } Q^{(k)}_{\beta} \mbox{ be square matrices and } d^{(k)} \mbox{ is a column vector indexed} \\ \mbox{expressions} & \mbox{by } S^{(k)} \mbox{ defined as follows:} \end{array}$ 

for performace

$$\begin{split} P_{ij}^{(k)} &\coloneqq P_{ij}, \quad \forall i, j \in S^{(k)}, \\ Q_{\beta}^{(k)} &\coloneqq [I_{2k-1} - \beta P^{(k)}]^{-1}, \\ d^{(k)} &\coloneqq [d(-k+1), \dots, d(k-1)]^{\mathsf{T}} \end{split}$$

Then,

$$L_{\beta}^{(k)} = \big\langle [Q_{\beta}^{(k)}]_0, d^{(k)} \big\rangle \quad \text{and} \quad M_{\beta}^{(k)} = \big\langle [Q_{\beta}^{(k)}]_0, \mathbf{1}_{2k-1} \big\rangle.$$



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#### Proof Standard Markov chain analysis.



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 $D_{\beta}^{(k)}$  and  $N_{\beta}^{(k)}$  can be computed using these expressions.



We found performance of a generic (threshold-based) strategy How does this lead to identifying an optimal strategy?

Use the idea of calibration from multi-armed bandits. Instead of finding the best strategy for a particular  $\lambda$ , we identify the set of  $\lambda$  that are optimal for a particular strategy.



$$\text{Some inequalities} \qquad L_{\beta}^{(k)} < L_{\beta}^{(k+1)}, \quad M_{\beta}^{(k)} < M_{\beta}^{(k+1)}, \quad D_{\beta}^{(k)} < D_{\beta}^{(k+1)}.$$

**Proof idea**  $\blacktriangleright$  For the first two inequalities, express  $P^{(k+1)}$  in terms in  $P^{(k)}$ .

• For the third inequality, define operator  $T^{(k+1)}$  as

$$[T^{(k+1)}D](e) = \begin{cases} \beta \sum_{n=-\infty}^{\infty} P_{0n}D(n), & |e| \ge k+1\\ (1-\beta)d(e) + \beta \sum_{n=-\infty}^{\infty} P_{en}D(n), & |e| < k+1 \end{cases}$$

• Define  $D^{(k,0)} = D^{(k)}$  and  $D^{(k,m+1)} = T^{(k+1)}D^{(k,m)}$ .

Show that  $D^{(k,m)}(e) > D^{(k)}(e)$  for all  $e \in A^{(m)}$ , where  $A^{(m)} \uparrow \mathbb{Z}$ .

 $\text{Some inequalities} \qquad L_{\beta}^{(k)} < L_{\beta}^{(k+1)}, \quad M_{\beta}^{(k)} < M_{\beta}^{(k+1)}, \quad D_{\beta}^{(k)} < D_{\beta}^{(k+1)}.$ 

Lagrangian cost



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Some inequalities  $L_{\beta}^{(k)} < L$ 

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Lagrangian cost

 $C_{\beta}^{(k)}(\boldsymbol{\lambda}) \coloneqq C(f^{(k)}, g^*; \boldsymbol{\lambda}) = D_{\beta}^{(k)} + \boldsymbol{\lambda} N_{\beta}^{(k)}$ 



Optimal performance

For all λ ∈ (λ<sub>β</sub><sup>(k)</sup>, λ<sub>β</sub><sup>(k+1)</sup>] the threshold strategy f<sup>(k+1)</sup> is optimal.
C<sup>\*</sup><sub>β</sub>(λ) = min<sub>k∈Z</sub> C<sup>(k)</sup><sub>β</sub> is piecewise linear, continuous, concave, and increasing function of λ.


# **Step 6:** Back to the constrained optimization problem

 $\label{eq:stategy} \begin{array}{ll} \mbox{Bernoulli} & \mbox{Let } \theta \in [0,1] \mbox{ and } f_1 \mbox{ and } f_2 \mbox{ be two stationary strategies.} \\ \mbox{The Bernoulli randomized strategy } (f_1,f_2,\theta) \mbox{ randomizes between } f_1 \mbox{ and } f_2 \mbox{ at each stage, choosing } f_1 \mbox{ with probability } \theta \mbox{ and } f_2 \mbox{ with probability } \\ (1-\theta). \end{array}$ 

**Simple rand.** A Bernoulli randomized strategy  $(f_1, f_2, \theta)$  is simple if the actions strategy prescribed by  $f_1$  and  $f_2$  differ only at one state.

Main result

t Define  $k_{\beta}^* = \sup\{k \in \mathbb{Z}_{\geq 0} : N_{\beta}^{(k)} \geq \alpha\}$  and let  $\theta$  be such that

$$\theta N_{\beta}^{(k_{\beta}^{*})} + (1-\theta) N_{\beta}^{(k_{\beta}^{*}+1)} = \alpha$$

Then, the Bernoulli simple randomized strategy  $(f^{(k_{\beta}^{*})}, f^{(k_{\beta}^{*}+1)}, \theta)$  is optimal for the constrained optimization problem for  $\beta \in (0, 1]$ .



**Sufficient** A (possibly randomized) strategy  $(f^{\circ}, g^{\circ})$  is optimal for a constrained **condition for** optimization problem with  $\beta \in (0, 1]$  is the following conditions hold: optimality (C1)  $N_{\beta}(f^{\circ}, g^{\circ}) = \alpha$ .



Sufficient

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Completely characterize the distortiontransmission function and the optimal strategy that achieves any point on that function.

$$P_{ij} = \begin{cases} p, & \text{if } |i-j| = 1; \\ 1-2p, & \text{if } i = j; \\ 0, & \text{otherwise,} \end{cases} \quad \text{where } p \in (0, \frac{1}{2}), \quad d(e) = |e|$$



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$$P_{ij} = \begin{cases} p, & \text{if } |i-j| = 1; \\ 1-2p, & \text{if } i = j; \\ 0, & \text{otherwise,} \end{cases} \quad \text{where } p \in (0, \frac{1}{2}), \quad d(e) = |e|$$

**Discounted cost** Let  $K_{\beta} = -2 - (1 - \beta)/\beta p$  and  $m_{\beta} = \cosh^{-1}(-K_{\beta}/2)$ .

$$D_{\beta}^{(k)} = \frac{\sinh(km_{\beta}) - k\sinh(m_{\beta})}{2\sinh^{2}(km_{\beta}/2)\sinh(m_{\beta})}$$
$$N_{\beta}^{(k)} = \frac{2\beta p\sinh^{2}(m_{\beta}/2)\cosh(km_{\beta})}{\sinh^{2}(km_{\beta}/2)} - (1 - \beta)$$

Average cost 
$$D_1^{(k)} = \frac{k^2 - 1}{3k}$$
 and  $N_1^{(k)} = \frac{2p}{k^2}$ 



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$$\lambda_{\beta}^{(\kappa)}$$
 can be computed in terms of  $D_{\beta}^{(\kappa)}$  and  $N_{\beta}^{(\kappa)}$ 

Average cost 
$$D_1^{(k)} = \frac{k^2 - 1}{3k}$$
 and  $N_1^{(k)} = \frac{2p}{k^2}$ 

$$\lambda_1^{(k)} = \frac{k(k+1)(k^2+k+1)}{6p(2k+1)}$$

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$$\begin{split} \mathsf{D}_{\beta}^{(k)} &= \frac{\sinh(km_{\beta}) - k \sinh(m_{\beta})}{2 \sinh^2(km_{\beta}/2) \sinh(m_{\beta})} \\ \mathsf{N}_{\beta}^{(k)} &= \frac{2\beta p \sinh^2(m_{\beta}/2) \cosh(km_{\beta})}{\sinh^2(km_{\beta}/2)} - (1 - \beta) \\ \\ \\ \mathsf{k}_{\beta}^* &= \sup\left\{ k \in \mathbb{Z}_{\geq 0} : \frac{\sinh^2(m_{\beta}/2) \cosh(km_{\beta})}{\sinh^2(km_{\beta}/2)} \geqslant \frac{1 + \alpha - \beta}{2\beta p} \right\} \end{split}$$
Average cost
$$\mathsf{D}_{1}^{(k)} &= \frac{k^2 - 1}{3k} \quad \text{and} \quad \mathsf{N}_{1}^{(k)} &= \frac{2p}{k^2} \end{split}$$

$$k_1^* = \left\lfloor \sqrt{\frac{2p}{\alpha}} \right\rfloor$$

Real-time transmission of Markov sources- (Aditya Mahajan)









# Summary and Conclusion

# Problem • Real-time transmission of a Markov source under constraints on the number of transmissions.

- Investigated both discounted and average cost infinite horizon setups.
- Modeled as a decentralized stochastic control problem with two decision maker.
- As long as the transmitter uses a symmetric threshold based strategy, the estimation strategy does not depend on the transmission strategy.
- The problem of find the "best response" transmitter is a centralized stochastic control problem.

#### **Main results** $\blacktriangleright$ Simple Bernoulli randomized strategies $(f^{(k^*)}, f^{(k^*+1)}, \theta)$ are optimal.

- $k^*$  and  $\theta$  can be computed easily.
- Characterize the distortion-transmission function

#### **References** > Chakravorty and Mahajan, Allerton 2014.

- Chakravorty and Mahajan, CDC 2014
- Full paper to be posted to arxiv soon.

